

## Model Checking – Endterm

Last name: \_\_\_\_\_

First name: \_\_\_\_\_

Student ID no.: \_\_\_\_\_

Signature: \_\_\_\_\_

- If you feel ill, let us know immediately.
- You have **100 minutes** to fill in all the required information and write down your solutions.
- Don't forget to **sign**.
- Write with a non-erasable **pen**, do not use red or green color.
- You are not allowed to use **auxiliary means**.
- You may answer in **English or German**.
- Please turn off your **cell phone**.
- Should you require additional **scrap paper**, please tell us.
- You can obtain **40 points** in the exam. You need **17 points** in total to pass (grade 4.0).
- Don't fill in the table below.
- Good luck!

Ex1	Ex2	Ex3	Ex4	Ex5	Ex6	$\Sigma$



**Exercise 1** LTL

2+2+2+2=8P

Consider the following LTL formulas over the set of atomic propositions  $AP = \{p, q\}$ :

$$\varphi_1 = \mathbf{G}(\mathbf{F}p \rightarrow q) \quad \varphi_2 = \mathbf{G}(q \mathcal{U} p) \quad \varphi_3 = \mathbf{G}(\mathbf{F}p \vee (\neg q \mathcal{U} \neg p))$$

- Is there a word satisfying  $\varphi_1$  but not  $\varphi_2$  ?  
If so, exhibit such a word and briefly explain why it does; else, justify briefly why it does not exist.
- Is there a word satisfying  $\varphi_2$  but not  $\varphi_1$  ?  
If so, exhibit such a word and briefly explain why it does; else, justify briefly why it does not exist.
- Is there a word satisfying all three formulas ?  
If so, exhibit such a word and briefly explain why it does; else, justify briefly why it does not exist.
- Give a Büchi automaton accepting exactly the words satisfying  $\varphi_2$ .

**Exercise 2** Spin

2+2+2= 6P

Consider the following Promela model:

```

1  byte g = 0;
2  chan c = [0] of { byte };
3
4  active proctype x() {
5      c!1;
6      do
7          :: c?0 -> g--; label1: c!1;
8      od
9  }
10
11 active proctype y() {
12     do
13         :: c?1 -> g++; c!0;
14         :: c?2 -> label2: g++;
15     od
16 }
17
18 ltl p1 { [] (x@label1 -> <> (g == 1)) }
19 ltl p2 { []<> (g == 0) }
```

Recall that  $x@label1$  (line 18) is **true** only in the states where the process  $x$  is at  $label1$ .

- Does the LTL formula  $p1$  at line 18 hold? Justify your answer.
- Write a process named  $z$  such that every execution of the code consisting of the processes  $x, y, z$  executes the statement at  $label2$  in line 14 *exactly* once. Use the following template:  
`active proctype z() { * Here comes your code * }`
- Does the LTL formula  $p2$  at line 19 hold after the process  $z$  is added? Justify your answer.

**Exercise 3** LTL to Büchi translation

2+3+3= 8P

Consider the LTL formula  $\phi$  over the set of atomic propositions  $AP = \{p, q, r\}$  (each state is marked with the atomic propositions it satisfies):

$$\phi = [(p \mathcal{U} q) \rightarrow (r \mathcal{R} q)] \mathcal{U} (\mathbf{X}(p \wedge q \wedge r))$$

- Put  $\neg\phi$  in negation normal form (NNF), i.e., give a formula  $\psi$  in NNF equivalent to  $\neg\phi$ .  
OBSERVE: you have to put  $\neg\phi$  in NNF, not  $\phi$ !
- Give a smallest consistent state of the automaton containing the following subformulas:

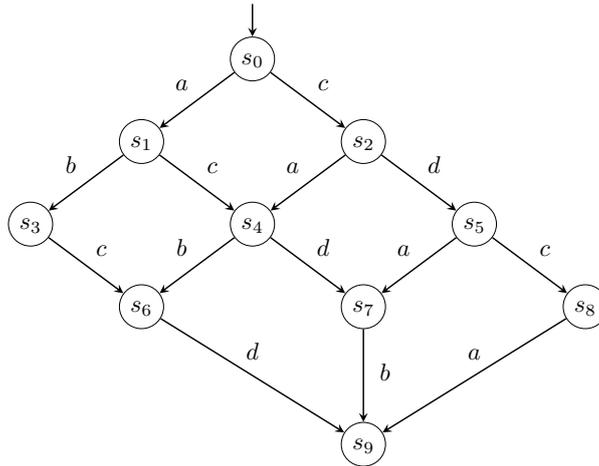
$$\psi, \mathbf{X}(\neg p \vee \neg q \vee \neg r), p, q, \neg r$$

- Give the target state and the label of a transition of the automaton having the state of (b) as source.

**Exercise 4 Partial Order Reduction**

**2+3=5P**

Consider a labelled Kripke structure  $\mathcal{K} = (S, A, \rightarrow, r, AP, \nu)$  where  $S$ ,  $A$ ,  $\rightarrow$ , and  $r$  are graphically represented as follows:

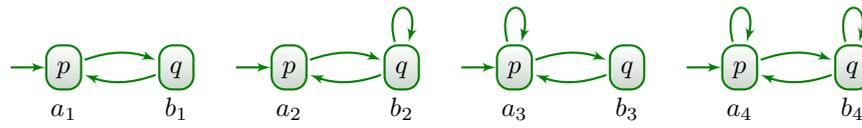


- (a) Give all pairs of actions that violate the diamond condition of independence relations. Justify your answer.
- (b) Assume that only the actions  $a$  and  $b$  are visible. Recall that  $red(s)$  denotes the set of actions  $a$  for which there is a transition  $s \xrightarrow{a} s'$  after the reduction. Can we have  $red(s_7) = \emptyset$  and satisfy conditions  $C_0$ - $C_3$ ? And  $red(s_1) = \{c\}$ ? And  $red(s_0) = \{c\}$ ? Justify your answers.

**Exercise 5 CTL**

**2+4= 6P**

Consider the following Kripke structure over the set of atomic propositions  $AP = \{p, q\}$ :



- (a) Give the set of states satisfying the CTL formulas:

$$\psi_1 = \text{AGEF}p \quad \psi_2 = \text{AFEG}p$$

- (b) Apply the fixpoint algorithm of the lecture to compute the set of states satisfying  $\text{AFEXEG}p$ . Describe briefly the intermediate steps.

**Exercise 6 BDDs**

**3+4= 7P**

- (a) Exhibit a BDD with 4 variables and 11 states (including the 0 state).
- (b) Prove that no BDD for a boolean function of arity 4 can have more than 11 states.